

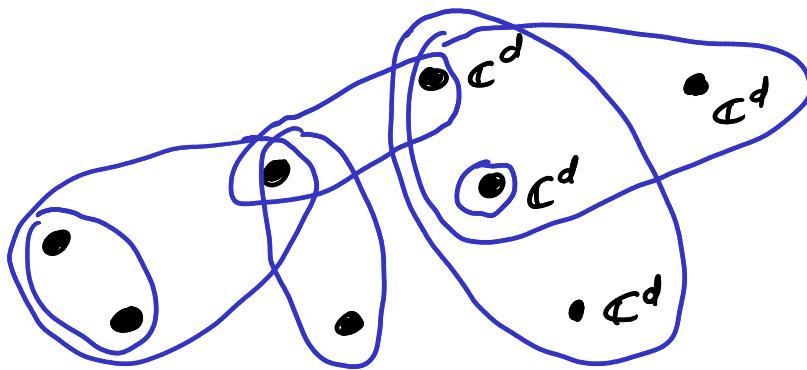
I. Hamiltonian Complexity

Notation & Terminology

- $\lambda_0(H)$ or $\lambda_{\min}(H) := \min_{|\phi\rangle} \langle \phi | H | \phi \rangle$
 "ground state energy"
- $H |\psi_0\rangle = \lambda_0 |\psi_0\rangle$ $|\psi_0\rangle$ = "ground state"
- $L_0 = \text{span} \{ |\psi\rangle : H|\psi\rangle = \lambda_{\min} |\psi\rangle \}$
 L_0 = "g.s. subspace"
- Note: $H \geq 0$, $\lambda_{\min}(H) = 0 \Rightarrow$ g.s. subspace = $\ker H$
- $\Delta(H) = \lambda_i - \lambda_0 = \min_{|\phi\rangle \perp L_0} \langle \phi | H | \phi \rangle - \lambda_0(H)$
 "spectral gap": same operator, different eigvals
- $\lambda_0(H_1) - \lambda_0(H_2)$
 "promise gap": different operators, same eigvals

Local Hamiltonians

Core object of study for this entire course.



Many-body quantum system:

- Multipartite Hilbert space $\mathcal{H} = (\mathbb{C}^d)^{\otimes n}$
(often $d=2$: n qubit system)
e.g. particles with spin- s ($d = \frac{s(s+1)}{2}$)

- Local interaction Hamiltonian:

$$h = h_S \otimes \mathbb{1}_{[n] \setminus S}, \quad S \subset [n]$$

i.e. h acts non-trivially only on subset S of the particles

We say that interaction h is k -local if $|S| = k$ (i.e. acts non-trivially on k particles).

Notation: often write " h_S " $\equiv h_S \otimes \mathbb{1}_{[n] \setminus S}$, i.e. subscript indicates indices of particles acted on, $\mathbb{1}$ on rest is implicit.

Def (k -local Hamiltonian)

$$H = \sum_i h^{(i)} \in \mathcal{B}((\mathbb{C}^d)^{\otimes n})$$

We say that H is a k -local Hamiltonian (or " H is k -local") if $\forall i$ $h^{(i)}$ is k -local.

I.e. k -local many-body Hamiltonian made of many interactions, each involving at most k particles.

Example: spins on line with nearest-neighbour interactions: $H = \sum_{i=1}^{n-1} h_{i,i+1}$

spins on lattice 
 $\text{---}, \sim : H = \sum_{\langle i,j \rangle} h_{ij}$
↗ neighbours

Note: In general, **no** underlying geometry
 \Rightarrow **no** requirement that local interactions are geometrically local.

Physics terminology: " k -local" = " k -body" but stuck with " k -local" now in Q.I.T.

Motivation:

Fundamental interaction n nature
 (e.g. electromagnetic, or "Coulomb", interaction)
 are not k-local — interactions decay
 as $\frac{1}{\text{poly}}$, not finite-range.

However, in many-body system electron clouds of neighbouring atoms shield atoms from fields of more distant particles (cf. Faraday cage).

At low-energies, many-body systems often well-approximated by k-local (often 2-local) Hamiltonians.

→ k-local H crop up throughout condensed matter physics.

The major topic of cond. mat.: phase transitions.

Quantum phase transitions (e.g. superconductivity, superfluidity) occur at zero or very low temperature characterised by abrupt change in ground state of system

→ quantum condensed matter ≈ study of ground states of many-body systems.

The "Local Hamiltonian problem" asks whether or not a local Hamiltonian has a low energy ground state.

Def (Local Hamiltonian Problem)

Input: k -local Hamiltonian H on n qudits with m local terms,
 "Promise" {where $\lambda_0(H) \leq \alpha$ or $\lambda_0(H) \geq \beta$
 (with $\beta - \alpha \geq \frac{1}{\text{poly}(n)}$)}

Output: YES if $\lambda_0(H) \leq \alpha$
 NO if $\lambda_0(H) \geq \beta$

Note that:

- Input is classical description of H , e.g. specify d^{2k} matrix elements for each of the m terms.
- Wlog $m \leq \binom{n}{k} = O(n^k) = \text{poly}(n)$ problem size (for fixed k, d) scales as $\text{poly}(n)$
 \longrightarrow take n to measure problem size

- (Strictly speaking, number of bits of data required to specify input also depends on precision of matrix entries. All our results will hold if matrix entries are restricted to $\text{poly}(n)$ digits of precision, so we ignore the input precision from now on.)
- Together with input, given promise that condition $\lambda_0(H) \leq \alpha$ or $\lambda_0(H) \geq \beta$, $\beta - \alpha = \frac{1}{\text{poly}(n)}$ holds for the input
 \rightarrow only required to solve problem under assumption condition holds.
- Conversely, to prove hardness results must show condition does hold for any H we construct.

Thm (Kitaev)

The Local Hamiltonian problem is QMA-hard

The first & most important result in Hamiltonian Complexity theory.

Has many interesting implications for QIT, CompSci, & especially physics (see later)

To prove QMA-hardness:

Transform any QMA prob. \rightarrow Local H. prob.

Recall Def. QMA:

\exists poly-sized q. circuit $U = U_T \cdot U_{T-1} \cdots U_2 \cdot U_1$, s.t.

YES: $\exists |w\rangle: \Pr(U|w\rangle \text{ outputs "1"}) \geq 2/3$

NO: $\forall |w\rangle: \Pr(U|w\rangle \text{ }) \leq 1/3$.

Idea:

"Encode" verifier circuit \rightarrow local Hamiltonian s.t. ground state encodes output of circuit.

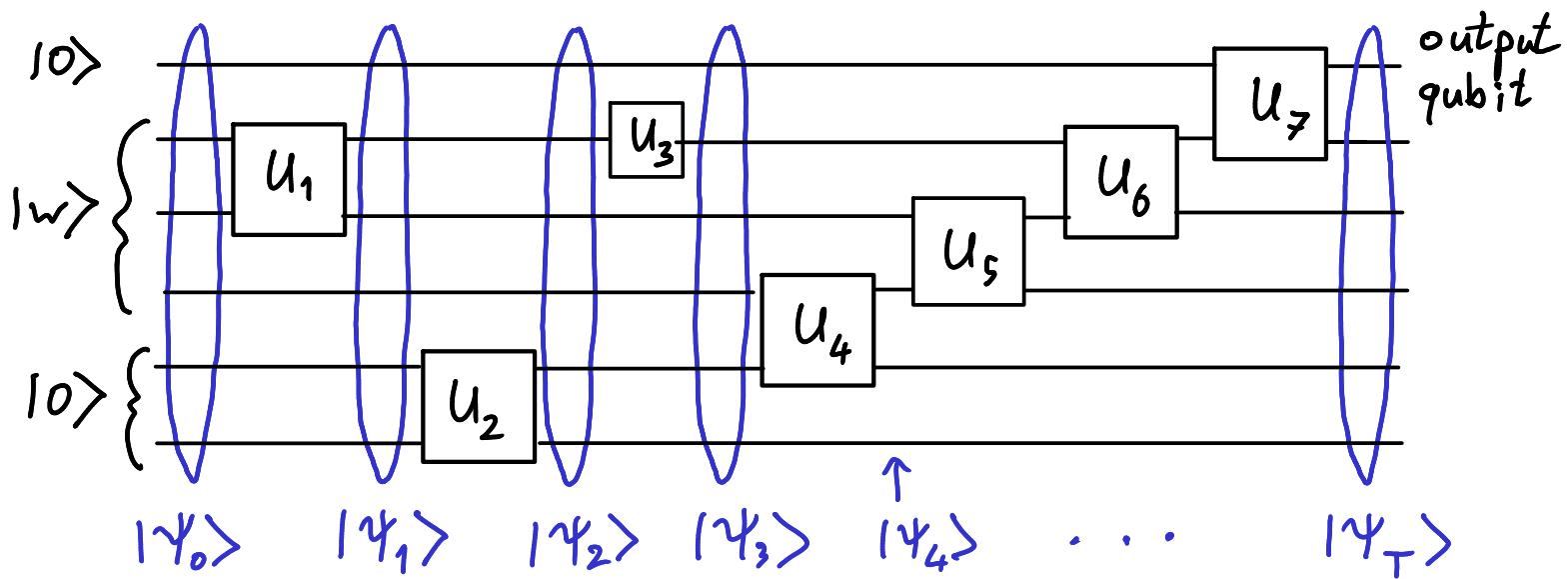
Add term that gives additional energy "penalty" if circuit outputs "0".

YES / NO \longleftrightarrow low-energy / high energy g.s.

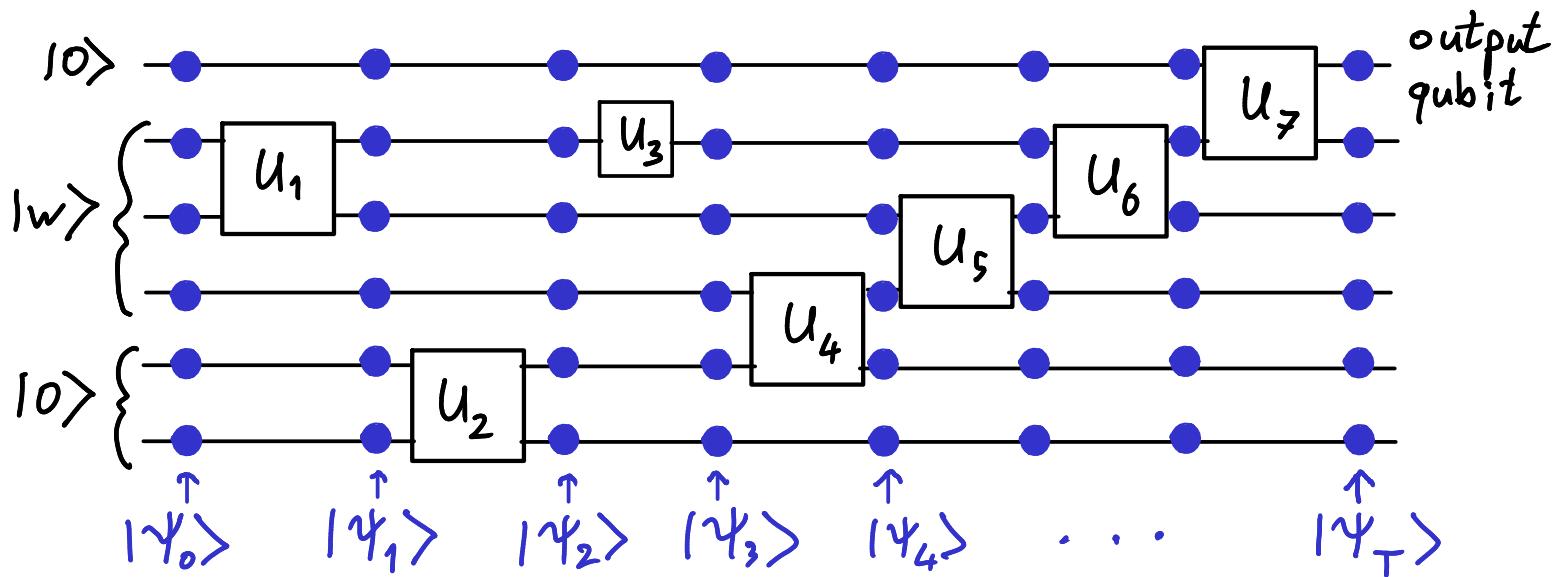
Before proving Thm, we first consider a naive approach that doesn't work, but teaches us something important about quantum many-body states and Hamiltonians.

Naive proof approach (doesn't work!)

Verifier circuit:

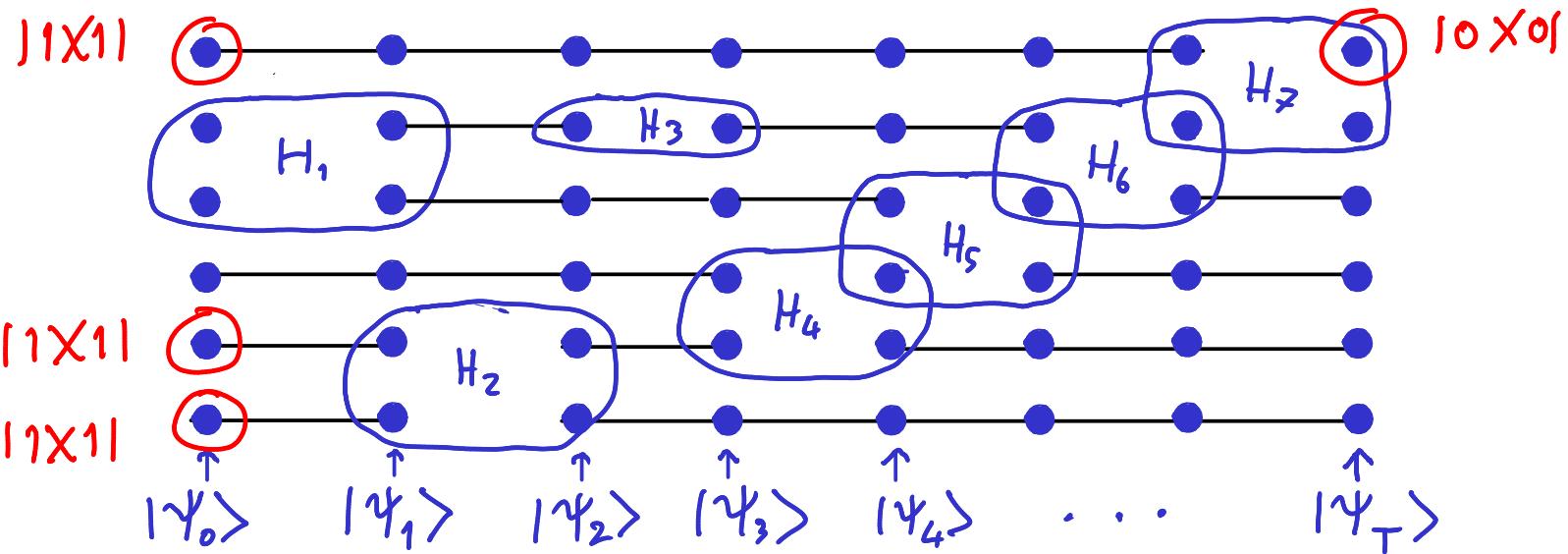


Imagine laying out circuit on lattice of qubits:



(9)

Now put local Hamiltonians wherever there's a gate (also need 2-local Hamiltonian terms for identity gates):



Can easily force qubits representing initial state of output & ancilla qubits to be $|0\rangle$ using 1-local term $\Pi^{(1)} = |1\rangle\langle 1|$,

But leave initial witness qubits state unconstrained

Similarly, can give additional energy penalty if qubit representing final state of output qubit is "0" using 1-local term $\Pi_{out}^{(0)} = |0\rangle\langle 0|$

If we could find Hamiltonian terms H_1, \dots, H_T s.t. ground states represent correct evolution of circuit, we'd be done:

YES instance:

\exists witness $|w\rangle$ s.t. $\Pr(\text{output } |1\rangle) \geq \frac{2}{3}$
 $\Rightarrow \exists$ state $|\Psi_0\rangle |\Psi_1\rangle \dots |\Psi_T\rangle$ representing correct evolution of circuit s.t.
only picks up energy $\leq \frac{1}{3}$ from $\Pi_{\text{out}}^{(0)}$
 \rightarrow g.s. energy $\leq \frac{1}{3}$

NO instance:

$\forall |w\rangle, \Pr(\text{output } |1\rangle) \leq \frac{1}{3}$
 $\Rightarrow \forall |\Psi\rangle = |\Psi_0\rangle |\Psi_1\rangle \dots |\Psi_T\rangle$ representing correct evolⁿ of circuit,
pick up energy $\geq \frac{2}{3}$ from $\Pi_{\text{out}}^{(0)}$
 $\forall |\Psi\rangle$ not representing correct evolⁿ,
pick up large energy penalty from H_1, \dots, H_T
 \rightarrow g.s. energy $\geq \frac{2}{3}$

Exercise

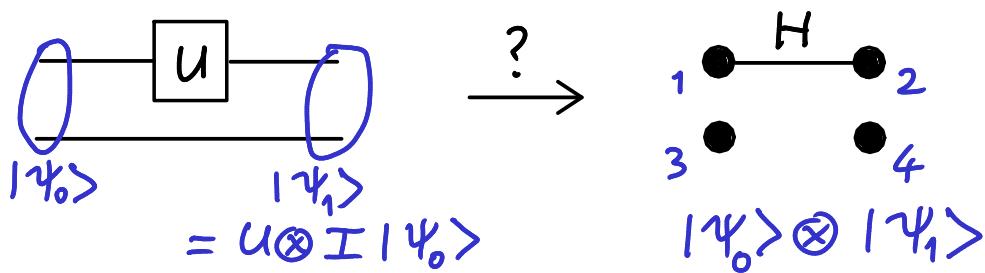
Make above argument rigorous.

If we could find Hamiltonian terms H_1, \dots, H_T s.t. ground states represent correct evolution of circuit, we'd be done.

Exercise

Use this approach to prove NP-hardness of Local Hamiltonian problem for classical Hamiltonians (diagonal in computational basis).

Unfortunately, such Hamiltonian terms cannot exist in quantum case.



Claim

$\exists U : \nexists H_{12}$ s.t. ground state subspace $L_0(H_{12} \otimes \mathbb{I}_{34}) = \text{span} \{ |\psi_{13}\rangle \otimes (U \otimes \mathbb{I}) |\psi_{24}\rangle \}$.

Proof

wLog can take $H \geq 0$, $\lambda_{\min}(H) = 0$ (substitute $H' = H + \lambda_{\min}(H) \cdot \mathbb{I}$, eigenvectors unchanged).

$$\rightarrow L_0(H \otimes \mathbb{I}) = \ker(H \otimes \mathbb{I}).$$

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Consider 2-qubit Hilbert space $\mathcal{H} = \mathbb{C}^2 \otimes \mathbb{C}^2$ with trivial gate $U = \mathbb{1}$, so

$$\text{span}\left\{\left|\psi\right\rangle_{13} \otimes (U \otimes \mathbb{1}) \left|\psi\right\rangle_{24}\right\} = \text{span}\left\{\left|\psi\right\rangle_{13} \left|\psi\right\rangle_{24}\right\}.$$

Let $|\phi\rangle = \frac{1}{\sqrt{2}} (|00\rangle + |11\rangle)$. entangled state! (ebit)

$|\phi\rangle |\phi\rangle \in A_0$, so

$$0 = \langle \phi |_{13} \langle \phi |_{24} (H_{12} \otimes \mathbb{1}_{34}) |\phi\rangle_{13} |\phi\rangle_{24}$$

$$= \text{tr} [H_{12} \otimes \mathbb{1}_{34} \cdot |\phi\rangle_{13} |\phi\rangle_{24} \langle \phi |_{13} \langle \phi |_{24}]$$

$$= \text{tr} [H_{12} \cdot \text{tr}_3 (|\phi\rangle_{13} \langle \phi|) \otimes \text{tr}_4 (|\phi\rangle_{24} \langle \phi|)]$$

$$= \text{tr} [H_{12} \cdot \mathbb{1}_1 \otimes \mathbb{1}_2]$$

$$= \text{tr } H.$$

$$H \geq 0, \text{tr } H = 0 \Rightarrow H = 0.$$

$$\mathcal{L}_0(H \otimes \mathbb{1}) = \ker(O \otimes \mathbb{1}) = \ker(O)$$

$$= \mathcal{H} \neq \text{span}\left\{\left|\psi\right\rangle \left|\psi\right\rangle\right\} \quad \square$$

Exercise

Show Proposition holds $\forall U$.