

## Local Clock Construction

Problem with binary clock is that cannot tell what time it by looking at only a few of the bits  $\rightarrow$  non-local H.

Idea: use unary-clock.

- Embedding  $\iota: \mathbb{C}^{T+1} \hookrightarrow (\mathbb{C}^2)^{\otimes T}$

$$\begin{aligned} \mathbb{C}^{T+1} &\xrightarrow{\quad} (\mathbb{C}^2)^{\otimes T} \\ |t\rangle &\xrightarrow{\quad} |1\rangle^{\otimes t} |0\rangle^{\otimes T-t} \\ &= \underbrace{|1, \dots, 1,}_{t} \underbrace{|0 \dots 0\rangle}_{T-t} \end{aligned}$$

(Label unary clock qubits  $1, \dots, T$ )

- Replace clock terms:

$$|0 \underset{cl}{\times} 0\rangle \rightarrow |0 \underset{1}{\times} 0\rangle$$

$$|T \underset{cl}{\times} T\rangle \rightarrow |1 \underset{T}{\times} 1\rangle$$

$$|t \underset{cl}{\times} t\rangle \rightarrow |1 \underset{t}{\times} 1\rangle \otimes |0 \underset{t+1}{\times} 0\rangle \quad t \neq 0, T$$

$$|t \underset{cl}{\times} t-1\rangle \rightarrow |1 \underset{t-1}{\times} 1\rangle \otimes |1 \underset{t}{\times} 0\rangle \otimes |0 \underset{t+1}{\times} 0\rangle \quad t \neq 0, T$$

$$|1 \underset{cl}{\times} 0\rangle \rightarrow |1 \underset{1}{\times} 0\rangle \otimes |0 \underset{2}{\times} 0\rangle$$

$$|T \underset{cl}{\times} T-1\rangle \rightarrow |1 \underset{T-1}{\times} 1\rangle \otimes |1 \underset{T}{\times} 0\rangle$$

(Partially) defines a linear mapping on operators

[Super-formally:

Pad mapping to define action on remaining

elems. It  $\times_{S1}$  & extend by linearity to

$$K : \mathcal{B}(\mathbb{C}^{T+1}) \longrightarrow \mathcal{B}((\mathbb{C}^2)^{\otimes T}).$$

Under this mapping,

$$H = H_{in} + H_{prop} + H_{out} \longrightarrow H'$$

where  $H'$  acts on unary clock subspace

$L = \text{span}\{|1\rangle^{\otimes t} |0\rangle^{\otimes T-1}\} \subset (\mathbb{C}^2)^{\otimes T}$  in just the same way as  $H$  acts on original clock  $\mathbb{C}^{T+1}$ .

[Super-formally:

$$\begin{array}{ccc} \mathbb{C}^{T+1} & \xrightarrow{H} & \mathbb{C}^{T+1} \\ \downarrow & & \downarrow \\ (\mathbb{C}^2)^{\otimes T} & \xrightarrow{K(H)} & (\mathbb{C}^2)^{\otimes T} \end{array} ]$$

Exercise: Prove this.

What goes wrong if we use

$$|t \times_{cl} t-1| \longrightarrow |1 \times_{t-1} 1| \otimes |1 \times_t 0| ?$$

Problem: What about "extra" part of Hilbert space? ( $\mathcal{L} = (\mathbb{C}^2)^{\otimes T} \setminus \cup(\mathbb{C}^{T+1})$ )

Add additional term to Hamiltonian:

$$H_{\text{stab}} = \mathbb{1} \otimes \sum_{t=1}^{T-1} |0\rangle\langle 0|_t \otimes |1\rangle\langle 1|_{t+1}$$

→ gives energy penalty to states not of form  $|1\dots, 1, 0, \dots, 0\rangle$   
(i.e. to all states in  $\mathcal{L}^\perp$ ).

### Claim

Same YES/NO-instance bounds for  $H' + H_{\text{stab}}$  as for  $H$ .

Proof is largely a matter of formalising the blindingly obvious!

Proof

Note :  $\ker H_{\text{stab}} = \mathcal{L} \supseteq \text{supp } H'$   
 $\text{supp } H_{\text{stab}} = \mathcal{L}^\perp$

$$H' = \begin{pmatrix} H & 0 \\ 0 & 0 \end{pmatrix}_{\mathcal{L}}, \quad H_{\text{stab}} = \begin{pmatrix} 0 & 0 \\ 0 & P \end{pmatrix}_{\mathcal{L}^\perp}$$

YES instance:

$$\lambda_{\min}(H' + H_{\text{stab}}) = \min [\lambda_{\min}(H'|_{\mathcal{L}}), \lambda_{\min}(H_{\text{stab}}|_{\mathcal{L}^\perp})]$$

by above Note

$$\leq \lambda_{\min}(H'|_{\mathcal{L}}) = \lambda_{\min}(H)$$

$$\leq \frac{\epsilon}{T+1} \text{ as before.}$$

NO instance:

Note: subspace  $\mathcal{L} \subset (\mathbb{C}^2)^{\otimes T}$  invariant under both  $H'$ ,  $H_{\text{stab}}$

(i.e.  $\forall |\psi\rangle \in \mathcal{L} : H'|\psi\rangle, H_{\text{stab}}|\psi\rangle \in \mathcal{L}$ ).

$\Rightarrow H' + H_{\text{stab}}$  decomposes as

$$H' = (H'|_{\mathcal{L}} + H_{\text{stab}}|_{\mathcal{L}}) \oplus (H'|_{\mathcal{L}^\perp} + H_{\text{stab}}|_{\mathcal{L}^\perp})$$

$$\begin{pmatrix} H & 0 \\ 0 & 0 \end{pmatrix} + \begin{pmatrix} 0 & 0 \\ 0 & P \end{pmatrix} = \begin{pmatrix} H & \\ & P \end{pmatrix}$$

Now

$$\bullet (H'|_{\mathcal{L}} + H_{\text{stab}}|_{\mathcal{L}}) = H'|_{\mathcal{L}} + 0 \stackrel{\cong}{=} H \geq \Omega\left(\frac{1-\sqrt{\varepsilon}}{T^3}\right).$$

from previously

$$\bullet (H'|_{\mathcal{L}^\perp} + H_{\text{stab}}|_{\mathcal{L}^\perp}) = 0 + H_{\text{stab}}|_{\mathcal{L}^\perp} = H_{\text{stab}}|_{\text{supp } H_{\text{stab}}} \\ = \left( \sum_{t=1}^T \mathbb{1} \otimes \Pi_t^{(0)} \otimes \Pi_{t+1}^{(1)} \right)|_{\text{supp}} \geq 1.$$

$$\therefore H' + H_{\text{stab}} \geq \Omega\left(\frac{1-\sqrt{\varepsilon}}{T^3}\right) \text{ as before.}$$

□

Overall Hamiltonian is:

$$H' + H_{\text{stab}} = H'_{\text{in}} + \sum_t H'_t + H'_{\text{out}} + H_{\text{stab}}$$

$$H'_{\text{in}} = (\Pi_1^{(1)} + \sum_{j \in A} \Pi_j^{(1)}) \otimes \begin{smallmatrix} 1 & 0 \\ 0 & 1 \end{smallmatrix} \quad \text{2-local}$$

$$H'_{\text{out}} = \Pi_1^{(0)} \otimes \begin{smallmatrix} 1 & 1 \\ 1 & 1 \end{smallmatrix} \quad \text{2-local}$$

$$H'_t = \frac{1}{2} \underbrace{1 \otimes (\begin{smallmatrix} 1 & 1 \\ t & t+1 \end{smallmatrix} \otimes \begin{smallmatrix} 1 & 0 \\ t+1 & t+1 \end{smallmatrix})}_{\text{2-local}} + \underbrace{\begin{smallmatrix} 1 & 1 \\ t+1 & t+2 \end{smallmatrix} \otimes \begin{smallmatrix} 1 & 0 \\ t+2 & t+2 \end{smallmatrix}}_{\text{2-local}}$$

$$- \frac{1}{2} \underbrace{U_{t+1} \otimes (\begin{smallmatrix} 1 & 1 \\ t & t+1 \end{smallmatrix} \otimes \begin{smallmatrix} 1 & 1 \\ t+1 & t+2 \end{smallmatrix} \otimes \begin{smallmatrix} 1 & 0 \\ t+2 & t+2 \end{smallmatrix})}_{\leq 5\text{-local}} - \text{h.c.}$$

(recall  $U_t$  is 1- or 2-qubit gate)

$$H_{\text{stab}} = \sum_{t=1}^{T-1} \begin{smallmatrix} 1 & 0 \\ t & t+1 \end{smallmatrix} \otimes \begin{smallmatrix} 1 & 1 \\ t+1 & t+1 \end{smallmatrix} \quad \text{2-local}$$

$\therefore$  Overall Hamiltonian  $H' + H_{\text{stab}}$  is 5-local.

Have proven Local Hamiltonian problem is QMA-hard for  $k$ -local Hamiltonians with  $k \geq 5$ !

Exercise: Prove Local Ham.  $\in$  QMA.

QED (kitaev's thm)