

# QMA-hardness of Local Hamiltonian

Recall:

Def (QMA + amplification)

Decision (YES/NO) problem, problem size  $n$ .

Problem  $\in$  QMA if

$\exists T = O(\text{poly}(n))$ ,  $\varepsilon = \Omega\left(\frac{1}{\text{poly}(n)}\right)$ ,

quantum circuit  $U = U_T \cdot U_{T-1} \cdots U_2 \cdot U_1$  s.t.

YES:

$\exists |w\rangle : \Pr(U|w\rangle \text{ outputs "1"}) \geq 1 - \varepsilon$

NO:

$\forall |w\rangle : \Pr(U|w\rangle \text{ .. } n) \leq \varepsilon$

Def (Local Hamiltonian Problem)

Input:  $k$ -Local Hamiltonian  $H$  on  
n qudits with m local terms,  
where  $\lambda_o(H) \leq \alpha$  or  $\lambda_o(H) \geq \beta$   
with  $\beta - \alpha \geq \frac{1}{\text{poly}(n)}$

Output: YES if  $\lambda_o(H) \leq \alpha$   
NO if  $\lambda_o(H) \geq \beta$

Thm (Kitaev)

The Local Hamiltonian problem is  
QMA-hard

We saw that entanglement defeats naive proof approach.

→ Need smarter way of encoding witness verification computation into Hamiltonian.

Idea (Feynman): encode evolution of computation in quantum superposition:

"Computational history state" or "history state":

$$|\Psi\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |\psi_t\rangle |t\rangle_{cl}$$

$\xrightarrow{\text{computation register } (\mathbb{C}^2)^{\otimes n}}$        $\xrightarrow{\text{clock register } \mathbb{C}^{T+1}}$

$|\psi_t\rangle$  = state after  $t$  steps (gates)

$|\psi_0\rangle$  = initial input state

$|\psi_T\rangle$  = final output state

Clock register takes states  $\in \mathbb{C}^{T+1}$ :  
 $|0\rangle, |1\rangle, |2\rangle, \dots, |T-1\rangle, |T\rangle$

Easy to write down (non-local) Hamiltonian with  $|\Psi\rangle$  as ground state:

$$H = H_{\text{in}} + H_{\text{prop}}$$

$$= (\mathbb{1} - |\psi_0 \times \psi_0|) \otimes I \otimes \mathbb{1}$$

ensures correct initial state

$$+ \sum_{t=0}^{T-1} \left( |\psi_t \times \psi_t| \otimes |t \times t| + |\psi_{t+1} \times \psi_{t+1}| \otimes |t+1 \times t+1| \right)$$

forces clock to "tick"

$$- |\psi_{t+1} \times \psi_t| \otimes |t+1 \times t|$$

forces 1 step of computation  
& increments clock

$$- |\psi_t \times \psi_{t+1}| \otimes |t \times t+1|)$$

necessary for Hermiticity;  
forces 1 step of computation  
backwards in time &  
decrements clock.

#### Exercise 4

Show that unique g.s. of  $H$  is  $|\Psi\rangle$ .

Challenge is to:

- construct local  $H$  for QMA verifier
  - prove YES instance  $\Rightarrow \lambda_0(H) \leq \alpha$
  - prove NO instance  $\Rightarrow \lambda_0(H) \geq \beta$
- $(\beta - \alpha \geq \frac{1}{\text{poly}(n)} \text{ promise gap})$

## Proof (Kitaev Thm.)

Hamiltonian (not fully local yet!)

$$H = H_{in} + H_{prop} + H_{out}$$

Initially, take computation register

$$\mathcal{H} = \underbrace{(\mathbb{C}^2)^{\otimes n}}_{\text{computation register}} \otimes \underbrace{\mathbb{C}^{T+1}}_{\text{clock register}} = \underbrace{\mathbb{C}^2 \otimes (\mathbb{C}^2)^{\otimes |W|}}_{\substack{\text{output qubit} \\ \uparrow}} \otimes \underbrace{(\mathbb{C}^2)^{\otimes |A|}}_{\substack{\text{witness register} \\ \uparrow}} \otimes \underbrace{\mathbb{C}^{T+1}}_{\substack{\text{ancillas A} \\ \uparrow \text{qubits } W}}$$

(will make clock local later).

$$H_{in} := \Pi_1^{(1)} \otimes I_{cl} \times 01 + \sum_{j \in A} \Pi_j^{(1)} \otimes I_{cl} \times 01$$

forces output qubit & ancillas to "initially" be in  $|0\rangle$  state

$$H_{prop} := \sum_{t=1}^{T-1} H_t, \text{ where}$$

$$H_t := \frac{1}{2} \mathbb{1} \otimes (I_t \times tI + I_{t+1} \times t+1I)$$

forces clock to "tick"

$$- \frac{1}{2} U_{t+1} \otimes I_{t+1} \times tI - \frac{1}{2} U_{t+1}^\dagger \otimes I_t \times t+1I$$

forward computation step      backwards step

$$H_{out} := \Pi_1^{(0)} \otimes I_T \times T I$$

gives energy "penalty" if output of QMA verifier computation is "0".

## Lemma

$$H_{\text{prop}} \sim \mathbb{1} \otimes E$$

(i.e.  $\exists$  unitary  $W$  s.t.  $W H_{\text{prop}} W^\dagger = \mathbb{1} \otimes E$ )

$$\text{where } E = \sum_t \frac{1}{2} (\langle t \rangle - \langle t+1 \rangle) (\langle t \rangle - \langle t+1 \rangle)$$

$$= \begin{pmatrix} \frac{1}{2} & -\frac{1}{2} & & & \\ -\frac{1}{2} & 1 & -\frac{1}{2} & & \\ & -\frac{1}{2} & 1 & \ddots & \\ & & \ddots & \ddots & \\ & & & 1 & -\frac{1}{2} \\ & & & & -\frac{1}{2} & \frac{1}{2} \end{pmatrix}$$

## Proof

Define unitary

$$W := \sum_{t=0}^T \prod_{i=1}^t U_i^\dagger \otimes |t\rangle\langle t| \quad (\text{where } \prod_{i=1}^0 U_i := \mathbb{1})$$

"undoes"  $t$  steps of computation

$$= \mathbb{1} \otimes |0\rangle\langle 0| + \sum_{t=1}^T (U_1^\dagger \cdot U_2^\dagger \cdots U_{t-1}^\dagger \cdot U_t^\dagger) \otimes |t\rangle\langle t|$$

$$W H_t W^\dagger$$

$$= (U_1^\dagger \cdots U_t^\dagger \otimes |t\rangle\langle t| + U_1^\dagger \cdots U_t^\dagger \cdot U_{t+1}^\dagger \otimes |t+1\rangle\langle t+1|)$$

$$\cdot \frac{1}{2} (\mathbb{1} \otimes (|t\rangle\langle t| + |t+1\rangle\langle t+1|))$$

$$- \frac{1}{2} U_{t+1}^\dagger \otimes |t+1\rangle\langle t+1| - \frac{1}{2} U_{t+1}^\dagger \otimes |t\rangle\langle t+1|)$$

$$\cdot (U_t \cdots U_1 \otimes |t\rangle\langle t| + U_{t+1} \cdots U_t \cdots U_1 \otimes |t+1\rangle\langle t+1|)$$

$$\begin{aligned}
&= \frac{1}{2} \left( (\cancel{U_1^+ \cdots U_t^+})(\cancel{U_t \cdots U_1}) \otimes |t\rangle \cancel{\times} |t\rangle \right) \\
&\quad + (\cancel{U_1^+ \cdots U_t^+ U_{t+1}^+})(\cancel{U_{t+1}} \cancel{U_t \cdots U_1}) \otimes |t+1\rangle \cancel{\times} |t+1\rangle \\
&\quad - (\cancel{U_1^+ \cdots U_t^+ U_{t+1}^+}) \cancel{U_{t+1}} (\cancel{U_t \cdots U_1}) \otimes |t+1\rangle \cancel{\times} |t\rangle \\
&\quad - (\cancel{U_1^+ \cdots U_t^+}) \cancel{U_{t+1}^+} (\cancel{U_{t+1}} \cancel{U_t \cdots U_1}) \otimes |t\rangle \cancel{\times} |t+1\rangle
\end{aligned}$$

$$= \mathbb{1} \otimes |\phi_t \times \phi_t| \quad \text{where} \quad |\phi_t\rangle = \frac{1}{\sqrt{2}} (|t\rangle - |t+1\rangle)$$

$$H_{prop} \sim W H_{prop} W^\dagger = \sum_t \mathbb{1} \otimes |\phi_t \times \phi_t| = \mathbb{1} \otimes E.$$

□

YES instance:

$\exists$  witness  $|w\rangle$  s.t.

$\Pr(\text{circuit outputs "0" on input } |w\rangle) \leq \varepsilon$  (Def. QMA)

Let  $|\Psi\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |\Psi_t\rangle |t\rangle_{cl}$

where  $|\Psi_t\rangle = U_t \cdot U_{t-1} \cdots U_2 \cdot U_1 |\Psi_0\rangle$

$$|\Psi_0\rangle = |0\rangle |w\rangle \underbrace{|0\rangle \cdots |0\rangle}_{\substack{\text{output qubit} \\ \text{witness}}} \underbrace{\cdots}_{\text{ancillas}}$$

Note:

- $|\Psi\rangle = \left( \sum_t U_t \cdots U_1 \otimes |t\rangle \langle t| \right) (|\Psi_0\rangle \otimes \underbrace{\frac{1}{\sqrt{T+1}} \sum_t |t\rangle}_{|q_0\rangle})$   
 $= W^\dagger |\Psi_0\rangle \otimes |\Psi_0\rangle.$
- $\Pr(\text{output "0" on input } |w\rangle) = \langle \Psi_T | \Pi_1^{(0)} | \Psi_T \rangle \leq \varepsilon.$

Want to show  $\lambda_0(H) \leq \alpha$ ,  $\alpha$  small  
→ show computational history state has low energy.

$$\begin{aligned}
\lambda_0(H) &= \min_{|\phi\rangle} \langle \phi | H | \phi \rangle \\
&\leq \langle \Psi | H | \Psi \rangle \\
&= \langle \Psi | H_{in} | \Psi \rangle + \langle \Psi | H_{prop} | \Psi \rangle + \langle \Psi | H_{out} | \Psi \rangle \\
&= \sim " \sim + \sum_t \langle \Psi | H_t | \Psi \rangle + \sim " \sim
\end{aligned}$$

$$\begin{aligned}
\langle \Psi | H_{in} | \Psi \rangle &= \frac{1}{T+1} \left( \sum_{t'=0}^T \langle \Psi_{t'} | \langle t' | \right) \left( (\Pi_1^{(1)} + \sum_{j \in \text{anc}} \Pi_j^{(1)}) \otimes |0\rangle \langle 0| \right) \cdot \\
&\quad \cdot \left( \sum_{t=0}^T |\Psi_t\rangle \langle t| \right) \\
&\propto \langle \Psi_0 | \langle 0 | \left( (\Pi_1^{(1)} + \sum_{j \in \text{anc}} \Pi_j^{(1)}) \otimes |0\rangle \langle 0| \right) |\Psi_0\rangle |0\rangle \\
&= \langle 0 | \langle w | \langle 0 |^{\otimes A} \cdot \Pi_1^{(1)} \cdot |0\rangle |w\rangle |0\rangle \\
&= \langle 0 | \cancel{\Pi_1^{(1)}} |0\rangle = 0,
\end{aligned}$$

$$\begin{aligned}
\langle \Psi | H_{prop} | \Psi \rangle &= \langle \Psi_0 | \otimes \langle \varphi_0 | W H_{prop} W^\dagger |\Psi_0\rangle \otimes |\varphi_0\rangle \\
&= \langle \Psi_0 | \otimes \langle \varphi_0 | \mathbb{1} \otimes E |\Psi_0\rangle \otimes |\varphi_0\rangle \quad (W H_{prop} W^\dagger = \mathbb{1} \otimes E) \\
&= \langle \varphi_0 | E |\varphi_0\rangle \\
&\propto \langle \varphi_0 | \sum_t (|t\rangle - |t+1\rangle) (\langle t | - \langle t+1 |) \sum_{t'} |t'\rangle \\
&= 0.
\end{aligned}$$

$$\begin{aligned}
& \langle \Psi | H_{out} | \Psi \rangle \\
&= \frac{1}{T+1} \left( \sum_{t=0}^T \langle \Psi_t | \langle t' | \right) (\Pi_1^{(o)} \otimes I_T \otimes I_T) \cdot \\
&\quad \cdot \left( \sum_{t=0}^T |\Psi_t\rangle |t\rangle \right) \\
&= \frac{1}{T+1} \langle \Psi_T | \Pi_1^{(o)} | \Psi_T \rangle \\
&= \frac{1}{T+1} \Pr(\text{output "0" on input } \omega) \quad (\text{from above}) \\
&\leq \frac{\epsilon}{T+1}
\end{aligned}$$

Thus

$$\begin{aligned}
\lambda_o(H) &\leq \cancel{\langle \Psi | H_{in} | \Psi \rangle}^{=0} + \sum_{t=1}^T \cancel{\langle \Psi | H_t | \Psi \rangle}^{=0} \\
&\quad + \langle \Psi | H_{out} | \Psi \rangle \\
&\leq \frac{\epsilon}{T+1} \quad \text{YES instances .}
\end{aligned}$$

NO instance:

$\forall |w\rangle :$

$\Pr(\text{circuit outputs "0" on input } |w\rangle)$

$$= \langle \gamma_T | \pi_i^{(0)} | \gamma_T \rangle$$

$$\geq 1 - \varepsilon \quad (\text{Def. QMA})$$

Need following:

Def (angle between subspaces)

Subspaces  $\mathcal{L}_1, \mathcal{L}_2$ .

$$\cos \theta(\mathcal{L}_1, \mathcal{L}_2) := \max_{\begin{array}{l} |\varphi_1\rangle \in \mathcal{L}_1 \\ |\varphi_2\rangle \in \mathcal{L}_2 \end{array}} |\langle \varphi_1 | \varphi_2 \rangle|$$

$$(\text{cf. } \vec{u} \cdot \vec{v} = \cos \theta, \quad |\vec{u}| = |\vec{v}| = 1)$$

Lemma (Kitaev)

If  $A, B \geq 0$  and

$\lambda_{\min}(A|_{\text{supp } A}), \lambda_{\min}(B|_{\text{supp } B}) \geq \mu$ , then

$$A + B \geq 2\mu \sin^2 \frac{\theta}{2}$$

where  $\theta = \theta(\ker A, \ker B)$ .

(Recall  $A \geq c \Leftrightarrow A - c\mathbb{1} \geq 0 \Leftrightarrow \lambda_{\min}(A) \geq c$ .)

## Proof

If  $\ker A \cap \ker B \neq \emptyset$ ,  $\theta = 0$  & bound is trivially true. So assume wlog  $\ker A \cap \ker B = \emptyset$ .

Let  $\Pi_A$  = projector onto  $\ker A$   
 $\Pi_B$  = " " "  $\ker B$

$$A \geq \mu(1 - \Pi_A), \quad B \geq \mu(1 - \Pi_A) \quad (\text{trivial})$$

→ sufficient to prove

$$(1 - \Pi_A) + (1 - \Pi_B) \geq 2 \sin^2 \frac{\theta}{2}$$

or

$$\Pi_A + \Pi_B \leq 1 + \cos \theta. \quad (*)$$

Let  $|\psi\rangle$  be any eigvect. of  $\Pi_A + \Pi_B$  with eigenval  $\lambda$  ( $> 0$  wlog; (\*) trivially true for  $\lambda = 0$ ):

$$\lambda |\psi\rangle = (\Pi_A + \Pi_B) |\psi\rangle = a |\varphi_A\rangle + b |\varphi_B\rangle$$

$$\text{where } \Pi_A |\psi\rangle =: a |\varphi_A\rangle, \quad \Pi_B |\psi\rangle =: b |\varphi_B\rangle$$

$$\& \text{ wlog } a, b \in \mathbb{R}^+ \quad (\text{absorb phases into } |\varphi_{A,B}\rangle)$$

$$\lambda = \langle \psi | (\Pi_A + \Pi_B) | \psi \rangle$$

$$= \langle \psi | \Pi_A \Pi_A | \psi \rangle + \langle \psi | \Pi_B \Pi_B | \psi \rangle$$

( $\Pi^2 = \Pi$  for projectors)

$$= a^2 + b^2$$

$$\begin{aligned}
\lambda^2 &= \langle \psi | (\pi_A + \pi_B)^2 | \psi \rangle \\
&= \underbrace{a^2 + b^2}_{\lambda''} + 2ab \operatorname{Re} \langle \varphi_A | \varphi_B \rangle \\
&\leq \lambda'' + 2ab |\operatorname{Re} \langle \varphi_A | \varphi_B \rangle| \quad a, b \in \mathbb{R}^+ \\
&\leq \lambda + (a^2 + b^2) |\operatorname{Re} \langle \varphi_A | \varphi_B \rangle| \\
&\quad \text{↑ } (a^2 + b^2 - 2ab = (a-b)^2 \geq 0 \Rightarrow 2ab \leq a^2 + b^2) \\
&= \lambda (1 + |\operatorname{Re} \langle \varphi_A | \varphi_B \rangle|) \\
&\leq \lambda \left( 1 + \max_{\substack{|\varphi_A\rangle \in \ker A \\ |\varphi_B\rangle \in \ker B}} |\langle \varphi_A | \varphi_B \rangle| \right) \\
&= \lambda (1 + \cos \theta)
\end{aligned}$$

$\therefore \lambda \leq 1 + \cos \theta$  which proves (\*).

□

$$\text{Let } A = H_{\text{in}} + H_{\text{out}}$$

$$B = H_{\text{prop}} = \sum_t H_t$$

in Lemma.

Need to bound  $\lambda_{\min}(A|_{\text{supp}A})$ ,  $\lambda_{\min}(B|_{\text{supp}B})$ ,  $\Theta(\ker A, \ker B)$

Claim  $\lambda_{\min}(A|_{\text{supp}A}) \geq 1$

smallest  
non-zero  
eigval

Proof:  $H_{\text{in}}, H_{\text{out}}$  commuting projectors

Claim  $\lambda_{\min}(B|_{\text{supp}B}) \geq -\Omega(T^{-2})$

Proof

Define unitary

$$W := \sum_t (U_1^\dagger \cdot U_2^\dagger \cdots U_{t-1}^\dagger \cdot U_t^\dagger) \otimes I_{t \times t}$$

(We used this before in proof of  $H \geq 0$ .)

$$B = H_{\text{prop}} \sim W H_{\text{prop}} W^\dagger = \mathbb{1} \otimes E$$

where  $E = \begin{pmatrix} \frac{1}{2} & -\frac{1}{2} & & & \\ -\frac{1}{2} & 1 & -\frac{1}{2} & & \\ & -\frac{1}{2} & 1 & \ddots & \\ & & \ddots & & \\ & & & 1 & -\frac{1}{2} \\ & & & -\frac{1}{2} & \frac{1}{2} \end{pmatrix}$  from before

$\therefore$  Eigvals of  $B$  = eigvals of  $E$ .

Eigvals & eigvects of  $E$  are:

$$\lambda_k = 1 - \cos q_k , \quad |\psi_k\rangle \propto \sum_{j=0}^T \cos\left(q_k(j + \frac{1}{2})\right) |j\rangle$$

where

$$q_k = \frac{\pi k}{T+1} , \quad k = 0, \dots, T$$

Exercise Verify this.

$$\lambda_0 = 0, \quad |\psi_0\rangle = \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |t\rangle .$$

Smallest non-zero eigenval

$$\lambda_1 = 1 - \cos\left(\frac{\pi}{T+1}\right) \geq -2(T^{-2}). \quad \square$$

(15)

Claim  $\sin^2 \frac{\theta}{2} = \Omega\left(\frac{1-\sqrt{\varepsilon}}{\tau+1}\right), \quad \theta \equiv \theta(\ker A, \ker B)$

Proof

$$A = H_{in} + H_{out}$$

$$= (\Pi_1^{(1)} + \sum_{j \in A} \Pi_j^{(1)}) \otimes |0\rangle\langle 0|_d + \Pi_1^{(co)} \otimes |T\rangle\langle T|$$

$$\ker A = \text{span} \left\{ |0\rangle, |\psi\rangle \underbrace{|0\rangle \dots |0\rangle}_{\text{ancillas}} \otimes |0\rangle_{cl}, \right.$$

$$|\psi\rangle \otimes |t\rangle_{cl} \quad 0 < t < T,$$

$$\left. |1\rangle, |\phi\rangle \otimes |T\rangle_{cl} \right\}$$

$$\therefore \Pi_A = |0\rangle\langle 0| \otimes \mathbb{1} \otimes (|0\rangle \dots |0\rangle) (\langle 0| \dots \langle 0|) \otimes |0\rangle\langle 0|_d$$

$$+ \sum_{t=1}^{T-1} \mathbb{1} \otimes |t\rangle\langle t| + |1\rangle\langle 1| \otimes \mathbb{1} \otimes |T\rangle\langle T|_d$$

$$W \Pi_A W^\dagger = \overbrace{|0\rangle\langle 0| \otimes \mathbb{1} \otimes (|0\rangle \dots |0\rangle) (\langle 0| \dots \langle 0|) \otimes |0\rangle\langle 0|_d}^{\Pi_I}$$

$$+ \sum_{t=1}^{T-1} \mathbb{1} \otimes |t\rangle\langle t|$$

$$+ \underbrace{U^\dagger (|1\rangle\langle 1| \otimes \mathbb{1}_{E_n \setminus \{1\}}) U}_{\Pi_F} \otimes |T\rangle\langle T|_d$$

where  $U := U_T U_{T-1} \dots U_2 U_1$ .

$$\begin{aligned}
\cos^2 \theta &= \max_{\substack{|\xi\rangle \in \ker A \\ |\eta\rangle \in \ker B}} |\langle \eta | \xi \rangle|^2 \quad \text{by Def.} \\
&= \max_{|\xi\rangle, |\eta\rangle} |\langle \eta | \Pi_A |\xi \rangle|^2 \quad \Pi_A = \text{proj. on } \ker A \\
&= \max_{|\eta\rangle \in \ker B} \langle \eta | \Pi_A | \eta \rangle \quad \text{Cauchy-Schwarz} \\
&\quad (\text{saturate } |\langle \eta | \Pi_A |\xi \rangle|^2 \leq \|\Pi_A |\eta\rangle\| \cdot \|\xi\| = \langle \eta | \Pi_A | \eta \rangle) \\
&= \max_{|\eta\rangle \in \ker B} \langle \eta | W^\dagger (W \Pi_A W^\dagger) W | \eta \rangle \\
&= \max_{|\eta'\rangle \in \ker WBW^\dagger} \langle \eta' | \left( \Pi_I \otimes \sum_{t=1}^{T-1} \mathbb{1} \otimes I_t \otimes \mathbb{1} + \Pi_F \otimes \sum_{t=1}^T \mathbb{1} \otimes I_t \otimes \mathbb{1} \right) |\eta'\rangle \\
&= \frac{1}{T+1} \max_{|\varphi\rangle} \sum_{t=0}^T \langle \varphi | \langle t | \left( \dots \right) \sum_{t'=0}^T |\varphi\rangle |t'\rangle \\
&\quad W B W^\dagger = \mathbb{1} \otimes E \Rightarrow |\eta\rangle = |\varphi\rangle \otimes \frac{1}{\sqrt{T+1}} \sum_{t=0}^T |t\rangle \\
&= \frac{1}{T+1} \max_{|\varphi\rangle} \left( \langle \varphi | (\Pi_I + \Pi_F) |\varphi\rangle + T-1 \right) \\
&= \frac{1}{T+1} \max_{|\varphi\rangle} \langle \varphi | 2\mathbb{1} - \Pi_I^\perp - \Pi_F^\perp |\varphi\rangle + \frac{T-1}{T+1} \\
&\leq \frac{2 - 2 \sin^2 \frac{\theta}{2}}{T+1} + \frac{T-1}{T+1} \quad \text{by Kitaev Lemma} \\
&= \frac{1 + \cos \vartheta}{T+1} + \frac{T-1}{T+1} \\
&\quad \text{where } \vartheta = \theta(\text{supp } \Pi_I, \text{supp } \Pi_F).
\end{aligned}$$

$$\cos^2 \vartheta = \max_{\substack{|\eta\rangle \in \text{supp } \pi_I \\ |\xi\rangle \in \text{supp } \pi_F}} |\langle \eta | \xi \rangle|^2$$

$$= \max_{|w\rangle, |\psi\rangle} \left| \underbrace{\langle 0 | \langle w | \langle 0 | \dots \langle 0 |}_{= |\Psi_0\rangle} \cdot (U^\dagger |1\rangle, |\psi\rangle) \right|^2$$

$$\pi_I = |0\rangle \otimes |0\rangle \otimes \dots \otimes |0\rangle (\langle 0 | \dots \langle 0 |)$$

$$\Rightarrow |\eta\rangle = |0\rangle, |w\rangle |0\rangle \dots |0\rangle$$

$$\pi_F = U^\dagger (|1\rangle \otimes |1\rangle \otimes \dots \otimes |1\rangle) U$$

$$\Rightarrow |\xi\rangle = |1\rangle, |\psi\rangle$$

$$= \max_{|w\rangle, |\psi\rangle} (\langle \Psi_0 | U^\dagger) |1\rangle \otimes |1\rangle \otimes \dots \otimes |1\rangle (U | \Psi_0 \rangle)$$

$$\leq \max_{|w\rangle} \langle \Psi_T | \pi_1^{(1)} |\Psi_T \rangle$$

$\leq \varepsilon.$  NO instance!

$$\therefore \cos^2 \theta \leq \frac{1 + \cos \vartheta}{T+1} + \frac{T-1}{T+1} \leq 1 - \frac{1 - \sqrt{\varepsilon}}{T+1}$$

$$\sin^2 \frac{\theta}{2} = \frac{1 - \sqrt{\cos^2 \theta}}{2} \geq \frac{1 - \sqrt{\varepsilon}}{4(T+1)}$$

Taylor expand  
 $\sqrt{1-x} \geq 1 - \frac{x}{2}$

Summary:

$$\lambda_{\min}(A|_{\text{supp } A}) \geq 1$$

$$\lambda_{\min}(B|_{\text{supp } B}) \geq \Omega(T^{-2})$$

$$\sin^2 \frac{\Theta}{2} \geq \frac{1 - \sqrt{\varepsilon}}{4(T+1)}$$

Kitaev's Lemma  $\Rightarrow$  for NO instance have

$$H = A + B \geq \Omega\left(\frac{1 - \sqrt{\varepsilon}}{T^3}\right)$$

i.e. have lower bound on  $\lambda_{\min}$ :

$$\lambda_{\min}(H) = \Omega\left(\frac{1 - \sqrt{\varepsilon}}{T^3}\right) \quad \text{as required.}$$

We have proven:

$$\text{YES} \Rightarrow \lambda_{\min}(H) \leq \frac{\varepsilon}{T+1} =: \alpha$$

$$\text{NO} \Rightarrow \lambda_{\min}(H) \geq \Omega\left(\frac{1 - \sqrt{\varepsilon}}{T^3}\right) =: \beta$$

$$\text{Taking } \varepsilon = \Omega\left(\frac{1}{T^2}\right) \Rightarrow \beta - \alpha = \Omega\left(\frac{1}{T^3}\right)$$

$\therefore$  Have proven reduction from QMA to Local Hamiltonian ...

... except that clock register has dimension  $T+1$

→  $H$  not on qudits with const. local dim.  $d$ .

Could split clock into  $\log T$  qubits

= binary clock ( $\mathbb{C}^T \cong (\mathbb{C}^2)^{\otimes \log T}$ )

→ local terms  $\log T$ -local, not  $k$ -local ( $k$  const).

(Example of general phenomenon: can always trade off locality against local dimension.)